

## X-bar syntax

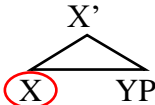
Universal inventory:

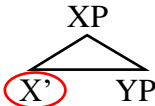
CAT is a set of possible categories.

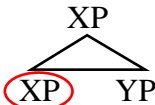
TRIANGLE is the set of possible triangles.

TRIANGLE = {COMP[X,Y], SPEC[X,Y], MOD[X,Y]: X, Y ∈ CAT}

where:

COMP[X,Y] = 

SPEC[X,Y] = 

MOD[X,Y] = 

(Encircled is the head)

BASE = {<α,l>, <α,r>: α ∈ TRIANGLE}

A grammar for language L is a pair <C<sub>L</sub>, S<sub>L</sub>>, where C<sub>L</sub> is a category assignment to the lexical items of L (a relation between LEX<sub>L</sub> and CAT) and S<sub>L</sub> is a syntax for L.

A syntax S<sub>L</sub> for language L is a subset of BASE.

We define, for α ∈ TRIANGLE, β ∈ {COMP, SPEC, MOD}

$\alpha =_L l$  iff  $\langle \alpha, l \rangle \in S_L$  and  $\langle \alpha, r \rangle \notin S_L$   
 $\alpha =_L r$  iff  $\langle \alpha, r \rangle \in S_L$  and  $\langle \alpha, l \rangle \notin S_L$   
 $\alpha =_L \perp$  iff  $\langle \alpha, l \rangle \in S_L$  and  $\langle \alpha, r \rangle \in S_L$   
 $\alpha =_L 0$  iff  $\langle \alpha, l \rangle \notin S_L$  and  $\langle \alpha, r \rangle \notin S_L$

$\beta_L = \{ \langle X, Y \rangle : \beta[X, Y] \neq_L 0 \}$

Grammar  $G_L$  for language  $L$  determines  $T_{G_L}$ , the tree set of  $G_L$ .

1.  $\alpha \in \text{LEX}_L$  and  $\langle \alpha, C \rangle \in C_L$  iff

$$\begin{array}{ccc} \begin{array}{c} C \\ | \\ \alpha \end{array} \in T_{G_L} & \begin{array}{c} C' \\ | \\ C \end{array} \in T_{G_L} & \begin{array}{c} CP \\ | \\ C' \end{array} \in T_{G_L} \end{array}$$

Plus pruning conditions: we can prune unary branches up to the highest node.

2.

$$\langle \text{COMP}[X, Y], l \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} X' \\ / \quad \backslash \\ X \quad YP \end{array} \in T_{G_L}$$

$$\langle \text{COMP}[X, Y], r \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} X' \\ / \quad \backslash \\ YP \quad X \end{array} \in T_{G_L}$$

$$\langle \text{SPEC}[X, Y], l \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} XP \\ / \quad \backslash \\ X' \quad YP \end{array} \in T_{G_L}$$

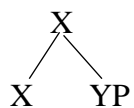
$$\langle \text{SPEC}[X, Y], r \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} XP \\ / \quad \backslash \\ YP \quad X' \end{array} \in T_{G_L}$$

$$\langle \text{MOD}[X, Y], l \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} XP \\ / \quad \backslash \\ XP \quad YP \end{array} \in T_{G_L}$$

$$\langle \text{MOD}[X, Y], r \rangle \in S_L \text{ iff} \quad = \quad \begin{array}{c} XP \\ / \quad \backslash \\ YP \quad XP \end{array} \in T_{G_L}$$

With this the specification

$\text{COMP}[X, Y] =_L l$  means that  $\langle \text{COMP}[X, Y], l \rangle \in S_L$  and  $\langle \text{COMP}[X, Y], r \rangle \notin S_L$ , and this means that the tree set  $T_{G_L}$  contains the tree, with the head  $X$  left of its complement  $YP$



With this the grammar constraints of, say, English are specified in a simple way:

$\langle V, D \rangle \in \text{COMP}_{\text{ENGLISH}}$

V takes DP complements

$\langle I, V \rangle \in \text{COMP}_{\text{ENGLISH}}$

I takes VP complements

$\langle I, D \rangle \in \text{SPEC}_{\text{ENGLISH}}$

I' takes DP specifiers

$\langle C, I \rangle \in \text{COMP}_{\text{ENGLISH}}$

C takes IP complements

$\langle D, N \rangle \in \text{COMP}_{\text{ENGLISH}}$

D takes NP complements

$\langle N, A \rangle \in \text{MOD}_{\text{ENGLISH}}$

NP takes AP modifiers

$\langle N, C \rangle \in \text{MOD}_{\text{ENGLISH}}$

NP takes CP modifiers

$\text{COMP}[V, D] =_{\text{ENGLISH}} l$

The head V is left of its DP complement

$\text{COMP}[I, V] =_{\text{ENGLISH}} l$

The head I is left of its VP complement

$\text{SPEC}[I, D] =_{\text{ENGLISH}} r$

The head I' is right of its DP specifier

$\text{COMP}[C, I] =_{\text{ENGLISH}} l$

The head C is left of its IP complement

$\text{COMP}[D, N] =_{\text{ENGLISH}} l$

The head D is left of its NP complement

$\text{MOD}[N, A] =_{\text{ENGLISH}} r$

The head NP is right of its AP modifier

$\text{MOD}[N, C] =_{\text{ENGLISH}} l$

The head NP is left of its CP modifier